



## ST10 UART recommendations

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### Introduction

This application note gives advice on using the ASC peripherals present in the ST10 and in particular, the asynchronous operations will be considered. A brief guide in the configuration of the peripheral and how to connect it to an RS232 interface or RS485 network is followed by an analysis of the sources of timing problems and how to minimize them to allow correct and error-free communications.

Please refer to the product user manual and datasheet for a full and detailed description of the functionality.

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# 1 ST10 ASC peripherals

Unlike the old microcontrollers of the ST10 family, two Asynchronous/Synchronous Serial (ASC) Interfaces are implemented on the latest ST10F27x and ST1025x, one as a standard peripheral and the other one mapped on the XBUS. The so-called standard module, instead, is present on all the devices of the ST10 family. The main differences between the two interfaces are limited to the programming model and interrupt management; in terms of general functionality and performance the two modules are equivalent.

The rest of the document refers only to the ST10F27x, recalling that all discussion regarding the standard module is still valid for all the other ST10 microcontrollers.

These peripherals provide serial communication between the ST10F27x and other microcontrollers, microprocessors or external peripherals. We refer to standard ASC interface as ASC0 and XBUS ASC interface as ASC1.

The ASC interfaces provide two kind of operations: Asynchronous and Synchronous. In synchronous mode, data is transmitted or received synchronously to a shift clock. In asynchronous mode, 8- or 9-bit data transfer, parity generation and a number of stop bits can be selected. Parity, framing and overrun error detection is provided to increase the reliability of data transfers. Transmission and reception of data is double-buffered.

For multiprocessor communication, a mechanism to distinguish address from data byte is included. Testing is supported by a loop-back option. A 13-bit baud rate generator provides the ASC with a separate serial clock signal.

## 1.1 Register linked with ASC modules

While ASC0 is enabled by default after the reset, ASC1 must be enabled by setting XPEN (bit 2 of SYSCON register) and bit 8 of XPERCON register. ASC0 is connected as an alternate function of bit 10 and bit 11 of PORT3 (P3.10/TXD0, P3.11/RXD0). ASC1 is connected as an alternate function of bit 6 and bit 7 of PORT8 (P8.6/TXD1 P8.7/RXD1). Obviously, RX pins must be set as input and TX pins as output, using the standard direction port register (DP3) for ASC0 and XS1PORT register for ASC1.

### 1.1.1 Control register

A control register controls the operating mode: S0CON for ASC0 and XS1CON for ASC1. The first one is bit addressable; the second one is not bit addressable. However, two additional registers (XS1CONSET and XS1CONCLR) are provided to set and clear single bits. Setting a bit in XS1CONSET register sets the corresponding bit in XS1CON register; clearing a bit has no effect. Likewise, setting a bit in XS1CONCLR register clears the corresponding bit in XS1CON register; clearing a bit has no effect.

SxCON (0 for ASC0 and 1 for ASC1)													Reset value: 0000h		
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SxR	SxLB	SxBRS	SxODD	-	SxOE	SxFE	SxPE	SxOEN	SxFEN	SxPEN	SxREN	SxSTP	SxM		
RW	RW	RW	RW		RW	RW	RW	RW	RW	RW	RW	RW		RW	

**Table 1. SxCON register description**

Bit	Function
SxM	ASCx mode control 0 0 0:8-bit data - synchronous operation 0 0 1:8-bit data - synchronous operation 0 1 0:Reserved. Do not use this combination 0 1 1:7-bit data + parity - asynchronous operation 1 0 0:9-bit data - asynchronous operation 1 0 1:8-bit data + wake up bit - asynchronous operation 1 1 0:Reserved. Do not use this combination 1 1 1:8-bit data + parity - asynchronous operation
SxSTP	Number of stop bit selection asynchronous operation 0:One stop bit 1:Two stop bit
SxREN	Receiver enable bit 0:Receiver disabled 1:Receiver enabled Reset by hardware after reception of byte in synchronous mode)
SxPEN	Parity check enable bit: asynchronous operation 0:Ignore parity 1:Check parity
SxFEN	Framing check enable bit - asynchronous operation 0:Ignore framing errors 1:Check framing errors
SxOEN	Overrun check enable bit 0:Ignore overrun errors 1:Check overrun errors
SxPE	Parity error flag Set by hardware on a parity error (SxPEN='1'). Must be reset by software.
SxFE	Framing error flag Set by hardware on a framing error (SxFEN='1'). Must be reset by software.
SxOE	Overrun error flag Set by hardware on an overrun error (SxOEN='1'). Must be reset by software.
SxODD	Parity selection bit 0:Even parity (parity bit set on odd number of '1's in data) 1:Odd parity (parity bit set on even number of '1's in data)
SxBRS	Baud rate selection bit 0:Divide clock by reload-value + constant (depending on mode) 1:Additionally reduce serial clock to 2/3rd
SxLB	Loopback mode enable bit 0:Standard transmit/receive mode 1:Loopback mode enabled
SxR	Baud rate generator run bit 0:Baud rate generator disabled (ASCx inactive) 1:Baud rate generator enabled

### 1.1.2 Transmit and receive buffer registers

Writing to the Transmit buffer register SxTBUF (via an instruction or a PEC data transfer) starts a transmission.

SxTBUF																Reset value: 0000h
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
-	-	-	-	-	-	-	Transmit Data Buffer									
-	-	-	-	-	-	-	RW									

SxRBUF																Reset value: 0xxxh
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
-	-	-	-	-	-	-	Received Data									
							R									

Only the number of data bit which is determined by the selected operating mode is actually transmitted. Bits written to positions 9 through 15 of register SxTBUF are always insignificant. After a transmission is completed, the transmit buffer register is cleared to 0000h.

Data transmission is double-buffered, so a new character may be written to the transmit buffer register before the transmission of the previous character is complete. This allows back-to-back transmission of characters without gaps.

Data reception is enabled by the Receiver Enable bit SxREN. After reception of a character is completed, the received data and, if provided by the selected operating mode, the received parity bit can be read from the (read-only) Receive Buffer register SxRBUF.

Bits in the upper half of SxRBUF which are not valid in the selected operating mode are read as zeros.

Data reception is double-buffered, so that reception of a second character may already begin before the previously received character is read out of the receive buffer register.

**The Loop-Back option** (selected by bit SxLB) allows the data currently being transmitted to be received simultaneously in the receive buffer.

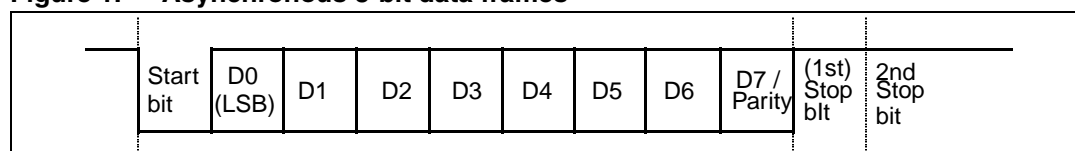
## 1.2 Asynchronous operation

In this document only Asynchronous Operation is considered; to use the ST10F27x ASC peripherals in synchronous mode, refer to the user manual. Asynchronous mode supports full-duplex communication, where both transmitter and receiver use the same data frame format and the same baud rate.

### 1.2.1 Asynchronous data frames

**8-bit data frames** either consist of 8 data bits D7...D0 (SxM = '001b') or of 7 data bits D6...D0 plus an automatically generated parity bit (SxM = '011b'). Parity may be odd or even, depending on bit SxODD in register SxCON. An even parity bit is set if the modulo-2-sum of the 7 data bits is '1', in which case an odd parity bit is cleared. Parity checking is enabled via bit SxPEN (always OFF in 8-bit data mode). If a wrong parity bit is received, the parity error flag SxPE is set along with the error interrupt request flag. The parity bit itself is stored in bit SxRBUF.7.

**Figure 1. Asynchronous 8-bit data frames**

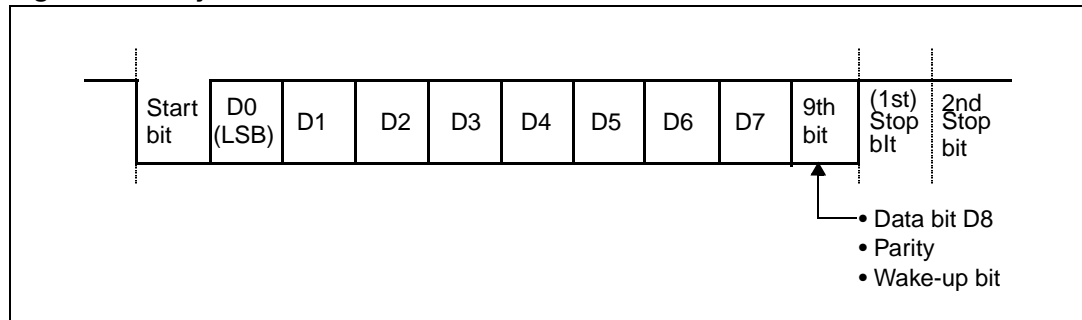


**9-bit data frames** either consist of 9 data bits D8...D0 (SxM = '100b'), or of 8 data bits D7...D0 plus an automatically generated parity bit (SxM = '111b') or of 8 data bits D7...D0 plus a wake-up bit (SxM = '101b'). Parity may be odd or even, depending on bit SxODD in register SxCON. An even parity bit is set, if the modulo-2-sum of the 8 data bits is '1', in which case an odd parity bit is cleared. Parity checking is enabled via bit SxPEN (always OFF in 9-bit data and wake-up mode). If a wrong parity bit is received, the parity error flag SxPE is set along with the error interrupt request flag. The parity bit itself is stored in bit 8 of SxRBUF.

In wake-up mode, received frames are only transferred to the receive buffer register if the 9th bit (the wake-up bit) is '1'. If this bit is '0', no receive interrupt request is activated and no data is transferred.

This feature may be used to control communication in a multiprocessor system. When the master processor wants to transmit a block of data to one of several slaves, it first sends out an address byte which identifies the target slave. An address byte differs from a data byte in that the additional 9th bit is a '1' for an address byte and a '0' for a data byte, so no slave is interrupted by a data byte. An address byte interrupts all slaves (operating in 8-bit data + wake-up bit mode), so each slave can examine the 8 LSBs of the received character (the address).

The addressed slave switches to 9-bit data mode (by clearing bit SxM.0), which enables it to also receive the data byte that will be coming (having the wake-up bit cleared). The slaves that were not being addressed remain in 8-bit data + wake-up bit mode, ignoring the following data byte.

**Figure 2. Asynchronous 9-bit data frames**

### 1.2.2 Asynchronous transmission

The ASCx module uses a modulo 16 counter clocked by  $f_{\text{CPU}} / 2$ . A transmission begins at the next overflow of this counter provided that SxR is set and data has been loaded into SxTBUF. The transmitted data frame consists of three basic elements:

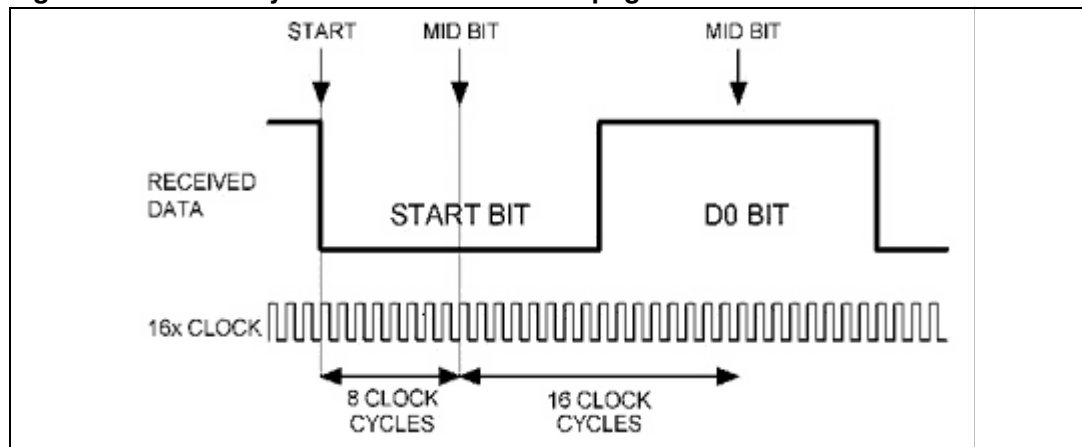
1. the start bit
2. the data field (8- or 9-bit, LSB first, including a parity bit, if selected)
3. the delimiter (1- or 2-stop bit)

Data transmission is double-buffered. When the transmitter is idle, the transmit data loaded into SxTBUF is immediately moved to the transmit shift register, thus freeing SxTBUF for the next data to be sent. This is indicated by the transmit buffer interrupt request flag being set. SxTBUF may now be loaded with the next data, while transmission of the previous one is still going on.

The transmit interrupt request flag is set before the last bit of a frame is transmitted, that is, before the first or the second stop bit is shifted out of the transmit shift register.

### 1.2.3 Asynchronous reception

A reception is initiated by a falling edge (1-to-0 transition) on pin RX, provided that bit SxR and SxREN are set. The receive data input pin is sampled at 16 times the rate of the selected baud rate. A majority decision of the 7th, 8th and 9th samples determines the effective bit value. This avoids erroneous results that may be caused by noise.

**Figure 3. ASCx asynchronous mode interrupt generation**

If the detected value is not a '0' when the start bit is sampled, the receive circuit is reset and waits for the next 1-to-0 transition at RX pin. If the start bit proves valid, the receive circuit continues sampling and shifts the incoming data frame into the receive shift register.

When the last stop bit is received, the content of the receive shift register is transferred to the receive data buffer register SxRBUF. Simultaneously, the receive interrupt request flag is set after the 9th sample in the last stop bit time slot (as programmed), regardless of whether or not a valid stop bit has been received. The receive circuit then waits for the next start bit (1-to-0 transition) at the receive data input pin.

Asynchronous reception is stopped by clearing bit SxREN. A currently received frame is completed including the generation of the receive interrupt request and an error interrupt request, if appropriate. A start bit that follows this frame is not recognized.

In wake-up mode, received frames are only transferred to the receive buffer register, if the 9th bit (the wake-up bit) is '1'. If this bit is '0', no receive interrupt request is activated and no data is transferred.

### 1.2.4 Hardware error detection

To improve the safety of serial data exchange, the serial channels provide an error interrupt request flag, which indicates the presence of an error, and three (selectable) error status flags in register SxCON, which indicate which error has been detected during reception. Upon completion of a reception, the error interrupt request flag is set simultaneously with the receive interrupt request flag, if one or more of the following conditions are met:

- If the framing error detection enable bit SxFEN is set and any of the expected stop bits are not high, the framing error flag SxFE is set, indicating that the error interrupt request is due to a framing error.
- If the parity error detection enable bit SxPEN is set in parity bit receive modes and the parity check on the received data bit proves false, the parity error flag SxPE is set, indicating that the error interrupt request is due to a parity error.
- If the overrun error detection enable bit SxOEN is set and the last character received was not read out of the receive buffer by software or PEC transfer at the time the reception of a new frame is complete, the overrun error flag SxOE is set indicating that the error interrupt request is due to an overrun error.

### 1.2.5 Baud rate generation

The serial channels have their own dedicated 13-bit baud rate generator. The baud rate generator is clocked by  $f_{CPU} / 2$ . The timer counts downwards and can be started or stopped through the Baud Rate Generator Run bit SxR in register SxCON. Each underflow of the timer provides one clock pulse to the serial channel. The timer is reloaded with the value stored in its 13-bit reload register each time it underflows. The resulting clock is again divided according to the operating mode and controlled by the Baud Rate Selection bit SxBRS. If SxBRS = '1', the clock signal is additionally divided to two-thirds of its frequency (see formulas and table). Thus, the baud rate is determined by the CPU clock, the reload value and the value of SxBRS.

S0BG (FEB4h / 5Ah)				SFR								Reset value: 0000h			
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
-	-	-	Baud Rate												
			RW												



For asynchronous operation, the baud rate generator provides a clock with 16 times the rate of the established baud rate. Every received bit is sampled at the 7th, 8th and 9th cycle of this clock. The baud rate for asynchronous operation of serial channels and the required reload value for a given baud rate can be determined by the following formulas:

$$B_{\text{Async}} = \frac{f_{\text{CPU}}}{16 \times [2 + (\text{SxBRS})] \times [(\text{SxBRL}) + 1]}$$

$$\text{SxBRL} = \left( \frac{f_{\text{CPU}}}{16 \times [2 + (\text{SxBRS})] \times B_{\text{Async}}} \right) - 1$$

(SxBRL) represents the content of the reload register, taken as unsigned 13-bit integer and (SxBRS) represents the value of bit SxBRS ('0' or '1'), taken as integer.

Using the above equation, the maximum baud rate can be calculated for any given clock speed. It follows a table of values for baud rate versus reload register value for SxBRS = 0 and SxBRS = 1.

**Table 2. Common baud rate settings and deviation error**

Baud rate	S1BRS = '0', $f_{\text{CPU}} = 64 \text{ MHz}$		S1BRS = '1', $f_{\text{CPU}} = 64 \text{ MHz}$	
	Deviation error	Reload value	Deviation error	Reload value
2 Mbaud	0%	0000 <sub>H</sub>	-	-
500 Kbaud	0%	0003 <sub>H</sub>	-	-
115.2 Kbaud	+2.12% / -3.55%	00010 <sub>H</sub> / 0011 <sub>H</sub>	+5.22% / -3.55%	000A <sub>H</sub> / 000B <sub>H</sub>
57.6 Kbaud	+2.12% / -0.79%	0021 <sub>H</sub> / 0022 <sub>H</sub>	+0.64% / -3.55%	0016 <sub>H</sub> / 0017 <sub>H</sub>
38.4 Kbaud	+0.16% / -1.73%	0033 <sub>H</sub> / 0034 <sub>H</sub>	+2.12% / -0.79%	0021 <sub>H</sub> / 0022 <sub>H</sub>
19.2 Kbaud	+0.16% / -0.79%	0067 <sub>H</sub> / 0068 <sub>H</sub>	+0.64% / -0.79%	0044 <sub>H</sub> / 0045 <sub>H</sub>
9600 baud	+0.16% / -0.32%	00CF <sub>H</sub> / 00D0 <sub>H</sub>	+0.64% / -0.08%	0089 <sub>H</sub> / 008A <sub>H</sub>
4800 baud	+0.16% / -0.08%	019F <sub>H</sub> / 01A0 <sub>H</sub>	+0.28% / -0.08%	0114 <sub>H</sub> / 0115 <sub>H</sub>
2400 baud	+0.04% / -0.08%	00CF <sub>H</sub> / 00D0 <sub>H</sub>	+0.10% / -0.08%	022A <sub>H</sub> / 022B <sub>H</sub>
1200 baud	+0.04% / -0.02%	0681 <sub>H</sub> / 0682 <sub>H</sub>	+0.01% / -0.08%	0456 <sub>H</sub> / 0457 <sub>H</sub>
600 baud	+0.01% / -0.02%	0D04 <sub>H</sub> / 0D05 <sub>H</sub>	+0.01% / -0.03%	08AD <sub>H</sub> / 08AE <sub>H</sub>
244 baud	-0.06%	1FFF <sub>H</sub>	+0.01% / -0.01%	1157 <sub>H</sub> / 1158 <sub>H</sub>

On the ST10 side, the baud rate is derived from the CPU frequency ( $f_{\text{CPU}}$ ). The  $f_{\text{CPU}}$  is a multiple (use of the internal PLL) of an external input clock which is generally different from the ones used in a PC to generate the UART clock. Therefore, there is a deviation error between the standard baud rate (reference being the PC) and the baud rate of the ST10 ASCx. This clock mismatch may lead to a reduction of tolerance as discussed in more detail in [Section 4.1: Arithmetic error in baud rate setting](#).

### 1.3 ASCx interrupt control

ASCx can generate four kinds of interrupt request as a result of four different events: Transmit, transmit buffer, receive and error. As already mentioned, the interrupt generation and management of ASC1 is quite different from the ASC0. In the ST10F27x all the peripherals connected on the XBUS share four IRQ lines and four registers (XIRxSEL, x = 0, 1, 2, 3) are provided in order to select the source of the XBUS interrupt. Whereas ASC1 provides four interrupt lines on the interrupt vectors, ASC0 uses its own four IRQ lines to manage the different events as illustrated in [Table 3](#):

**Table 3. ASCx interrupt lines**

Interrupt vectors	ASC1 interrupt lines	ASC0 interrupt lines
Receive	XP0INT, XP1INT, XP2INT	S0RINT
Transmit		S0TINT
Transmit Buffer		S0TBINT
Error	XP3INT	S0EINT

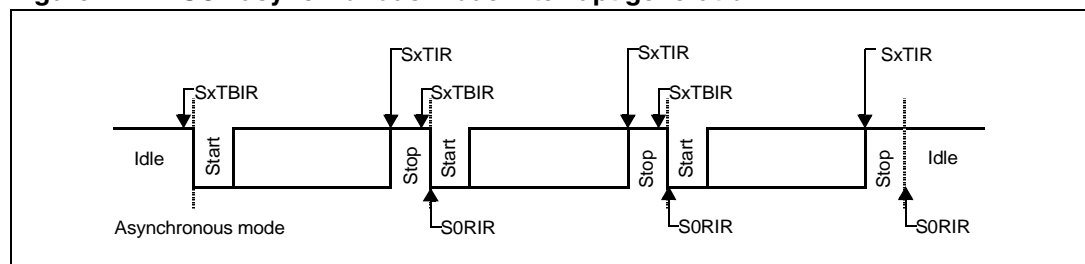
In addition, the registers used by ASC0 are bit addressable whereas XBUS registers are not. Other registers are provided to emulate the bit-addressability, as already explained in the description of XS1CONSET and XS1CONCLR (see [Section 1.1.1: Control register](#)). Please refer to the user manual for more details on interrupt generation.

An interrupt error request flag is set if one of the events described in [Section 1.2.4: Hardware error detection](#) occurs.

For normal operation (other than the error interrupt), the ASCx provides three interrupt requests to control data exchange via this serial channel:

- SxTBIR is activated when data is moved from SxTBUF to the transmit shift register.
- SxTIR is activated before the last bit of an asynchronous frame is transmitted or after the last bit of a synchronous frame has been transmitted.
- SxRIR is activated when the received frame is moved to SxRBUF.

**Figure 4. ASCx asynchronous mode interrupt generation**



## 2 Connecting an ASC to an RS232 interface

### 2.1 Introduction

The RS232 specification dates back to 1962, when it was first released by the EIA (Electronic Industries Association). The specification has changed over time, incorporating higher data rates and closing the compatibility gaps between TIA (Telecommunication Industry Association) and international (ITU, ISO) standards. The current version of the RS232RS232 specification is EIA/TIA-232-F, dated October 1997.

The RS232 link was initially intended to support modem and printer applications on IBM PCs, however, it now enables a variety of peripherals to communicate with PCs. The RS232 standard was defined as a single-ended standard for increasing serial-communication distances at low baud rates (< 20 Kbit/s).

### 2.2 Electrical characteristics

**Table 4. RS232 major electrical characteristics**

Parameter	Conditions	Min.	Max.	Unit
Driver output voltage open circuit			25	V
Driver output voltage loaded	$3k\Omega < R_L < 7k\Omega$	$\pm 5$	$\pm 15$	
Slew rate		4	30	V/ $\mu$ s
Maximum load capacitance			2500	pF
Receiver input resistance		3	7	k $\Omega$
Receiver input threshold				
Output = Mark (Logic 1)		-3		V
Output = Space (Logic 0)			3	

The major electrical characteristics of the RS232 standard are summarized in [Table 4](#) above. Typical RS232 transmissions seldom exceed 30 meters (often only a few meters) for two reasons. Firstly, the difference between transmitted levels ( $\pm 5V$ ) and receive levels ( $\pm 3V$ ) allows only 2V of common-mode rejection. Secondly, the distributed capacitance of a longer cable can degrade slew rates by exceeding the maximum specified load (2500pF). Because the RS232 was designed as a point-to-point rather than a multidrop interface, its drivers are specified for single loads from 3k to 7k ohm.

A typical RS232 signal fluctuates between positive and negative. Although the RS232 data is inverted, an overall translation from TTL/CMOS levels of the ST10F27x to RS232 as well as from RS232 to the TTL/CMOS levels of the ST10F27x is needed. A very common level translator is the ST232A, a device which includes a charge pump to generate  $\pm 12V$  from a single 5V supply, as well as two receivers and two transmitters in a single package. Please refer to the product datasheet for the full device description.

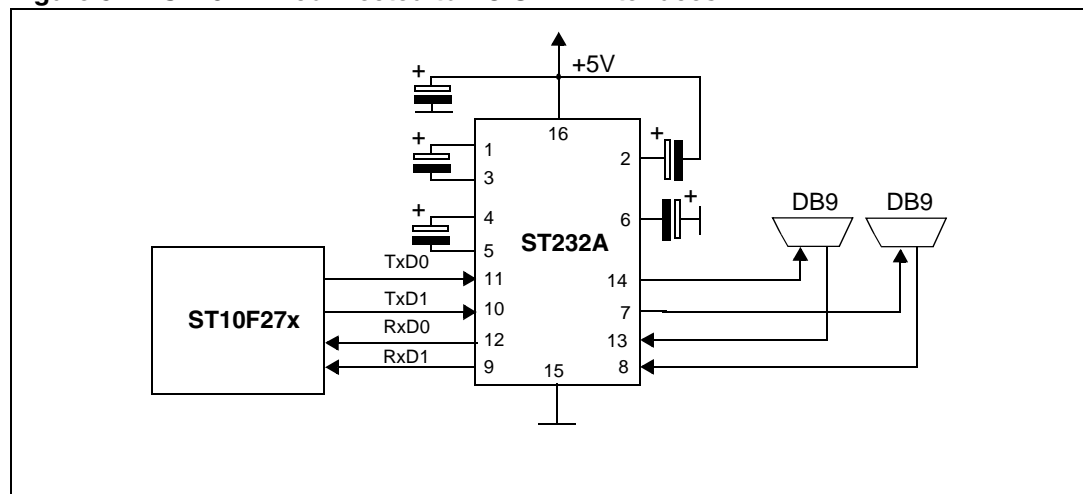
## 2.3 Connecting ST10F27x to a PC UART RS232 interface

A typical application involves a connection between a microcontroller and a personal computer. A PC UART interface uses a DB9 connector which contains more than only two signals and the ground. A certain number of other signals (such as RTS, DTR) are requested to set up a protocol between the PC and the peripheral (that is, modem). These signals are not necessarily required and often the serial transfers are performed using only Rx and Tx signals.

*Figure 5* below shows an example of a connection between two ASCx peripherals and two RS232 interfaces through an ST232A. The value of the capacitors needed for the charge pump circuitry is 0.1µF.

PC RS232 interfaces typically use baud rates of 9600...115200 bit/s. Since the clock sources are different, there is a deviation error between the standard baud rate and the baud rate of ST10F27x ASCx. This clock mismatch may lead to a reduction of tolerance as is discussed in more detail in [Section 4.1: Arithmetic error in baud rate setting](#).

**Figure 5. ST10F27x connected to PC UART interfaces**



## 3 Connecting ASC to an RS485 network

### 3.1 Introduction

An RS232 interface connects one DTE (data terminal equipment) to one DCE (data communication equipment) at a maximum speed of 20 Kbit/s with a maximum cable length of 30 meters. This was sufficient early on when almost all computer equipment was connected by modems but, soon after, interfaces capable of one or more of the following were in demand:

- Connecting DTE's directly without the need for modems
- Connecting several DTE's in a network structure
- Communicating over longer distances
- Communicating at faster speeds

RS485 is the most versatile communication standard in the standard series defined by the EIA, as it performs well on all four points. Thus RS485 is currently a widely-used communication interface in data acquisition and control applications where multiple nodes communicate with each other.

### 3.2 RS485 versus RS232

One of the main problems with RS232 is the lack of immunity against noise on the signal lines. The transmitter and receiver compare the voltages of the data and handshake lines with one common zero line. Shifts in the ground level can have disastrous effects. Therefore, the trigger level of the RS232 interface is set relatively high at  $\pm 3V$  ([Table 3](#)). Noise is easily picked up and limits both the maximum distance and communication speed. On the contrary, a common zero as a signal reference does not exist with RS485. A difference of several volts in the ground level of the RS485 transmitter and receiver does not cause any problems. The RS485 signals are floating and each signal is transmitted over a B line and an A line. The RS485 receiver compares the **voltage difference** between both lines, instead of the **absolute voltage level** on a signal line. This works well and prevents the existence of ground loops, a common source of communication problems. The best results are achieved if the B and A lines are twisted.

When the cable is twisted, we see that in some parts of the signal lines the direction of the noise current is the opposite from the current in other parts of the cable. Consequently, the resulting noise current is many factors lower than with an ordinary straight cable. Shielding tries to isolate the signal lines from hostile magnetic fields. Twisted pairs in RS485 communication, however, add immunity, which is a much better way to fight noise. Differential signals and twisting allow RS485 to communicate over much longer distances than achievable with RS232. With RS485, communication distances of 1200 meters are possible.

Differential signal lines also allow higher bit rates than possible with non-differential connections. Therefore, RS485 can overcome the practical communication speed limit of RS232. Currently, RS485 drivers are produced that can attain a bit rate of 35 Mbit/s.

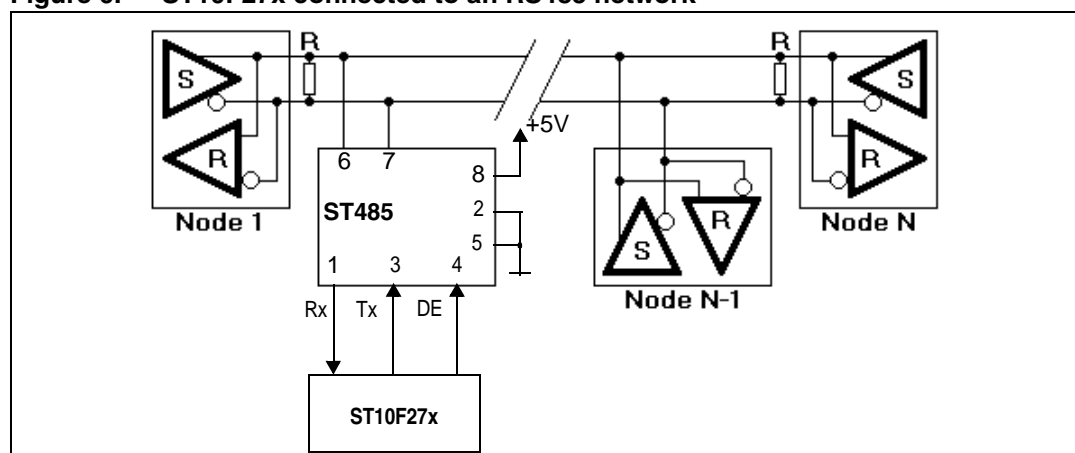
**Table 5. Comparison between RS232 and RS485**

Parameter	RS232	RS485
Differential	No	Yes
Maximum number of drivers	1	32
Maximum number of receivers		
Modes of operation	Full/Half Duplex	Half Duplex
Network topology	Point-to-Point	Multipoint
Maximum distance	30m	1200m
Maximum speed	20 Kbit/s	35 Mbit/s
Maximum slew rate	30V/ $\mu$ s	-
Receiver input resistance	3..7k $\Omega$	>12k $\Omega$
Driver load impedance		54 $\Omega$
Receiver input sensitivity	$\pm 3$ V	$\pm 200$ mV
Receiver input range	$\pm 15$ V	-7V..+12V
Maximum driver output voltage	$\pm 25$ V	
Minimum driver output voltage	$\pm 5$ V	$\pm 1.5$ V

### 3.3 Connecting ST10F27x to an RS485 network

[Figure 6 on page 14](#) shows an example of connection between the ST10F27x and an RS485 network through an RS485 transceiver (ST485). Please refer to the component datasheet for detailed information. Since each RS485 system must have a driver that can be disconnected from the transmission line when it is not transmitting, a generic ST10F27x GPIO pin, in addition to the ASCx standard pins TxDx and RxDx, can be connected to the ST485 DE (Driver Enable) pin to provide this functionality.

The problematic regarding data rate in relation with the cable characteristics and the correct termination of the bus are beyond the purpose of this application note (please refer to EIA/TIA-485 specifications).

**Figure 6. ST10F27x connected to an RS485 network**

## 4 UART communication and ST10 clock generation

The timing of a UART transmission is influenced by various factors. The baud rate of a UART is usually configured by integer values – the arithmetic rounding error leads to baud rate deviations (see [Table 2: Common baud rate settings and deviation error](#)). The architecture of ST10F27x UARTs furthermore leads to intrinsic delays at the sending UART. Moreover, the qualities of the clocks of the communication partners must be considered. Clock drift and jitter influence the baud rate and the instant of communication start. The network itself introduces some issues due to the finite channel bandwidth that reduces the useful time interval to correctly sample the level of the bit. This section quantifies these timing deviations separately, goes on to calculate the maximum baud rate deviation and provides some examples at the end.

### 4.1 Arithmetic error in baud rate setting

The baud rate depends on the frequency of the CPU ( $f_{CPU}$ ). The baud rate is set by choosing a value SxBRL (ASCx Baud Rate Setting) as follows:

$$SxBRL = \frac{f_{CPU}}{16 \times [2 + SxBRS] \times \text{baud}} - 1$$

Since SxBRL must be an integer, the value that rounds up the desiderated baud rate is:

$$SxBRL_{up} = \left\lceil \frac{f_{CPU}}{16 \times [2 + SxBRS] \times \text{baud}} \right\rceil - 1$$

while the value that rounds down the desiderated baud rate is:

$$SxBRL_{down} = \left\lfloor \frac{f_{CPU}}{16 \times [2 + SxBRS] \times \text{baud}} \right\rfloor$$

So the maximum error that is committed is:

$$\frac{1(SxBRL_{down} - SxBRL_{up})}{2} = \frac{8 \times [2 + SxBRS] \times \text{baud}}{f_{CPU}}$$

[Table 2: Common baud rate settings and deviation error on page 9](#) reports the arithmetic error committed when one of the standard baud rates is set, assuming the CPU clock to be 64 MHz. The last formula clearly shows that:

- the deviation error is directly proportional to the desiderated baud rate;
- choosing appropriately SxBRS, the error is always less than 2.12%.

## 4.2 Send/receive jitter problem

At initialization of the ASCx, the baud rate generator is started. Running at a frequency corresponding to the configured baud rate, the baud rate generator periodically generates ticks which are possible start events of a message bit. When the ASCx receives a transmit signal, it starts the transmission at the next tick from the baud rate generator. Thus, depending on the internal state of the baud rate generator, the transmission of a message may be delayed up to the interval time of two subsequent baud rate generator ticks. The state of the baud rate generator cannot be read, leading to an indeterministic send delay jitter:

$$0 \leq t_{\text{jitter}} \leq \frac{1}{\text{baud}}$$

Subsequent transmissions may not be affected by this send jitter, however if transmissions can start at particular instants, the send jitter jeopardizes the achievable accuracy for this send instant.

A similar situation exists for reception: A start condition is detected at the first tick from the clock that feeds the module after the falling edge, leading to a indeterministic receive delay jitter:

$$0 \leq t_{\text{jitter}} \leq \frac{1}{16 \cdot \text{baud}}$$

## 4.3 Clock issues

Usually components performing a UART communication are set apart from each other and are clocked by different clock sources. These clock sources are not absolutely free from errors. Finite accuracy, drift and jitter are common problems of clock generation.

In general, a standard, low cost, ceramic resonator with  $\pm 0.5\%$  accuracy and a further  $\pm 0.5\%$  drift over temperature and life can be used for the clock source at both ends of the link. On the other hand, a standard crystal oscillator usually provides 20ppm frequency tolerance and 100ppm frequency stability.

Another issue related to the clock is the use of a PLL (Phase Locked Loop). PLL is used in the ST10F27x to achieve higher internal clock frequencies but it introduces a small transient phase shift called jitter. In the ST10F27x the single period jitter is, in the worst case,  $\pm 500\text{ps}$  whereas the long term jitter can be maximum  $\pm 4.5\text{ns}$ . For more details regarding PLL jitter, please refer to the product datasheet.



## 4.4 Communication channel effects

In reception, the receiver must sample ideally in the middle of the bit; it is not desirable to sample close to the bit transition point. The main reason for this is the finite (and typically slow) transmission rise and fall times are made even slower if overly capacitive cabling is used. A long bus incurs high attenuation, which reduces noise margin, making it more important to sample when the bit level has settled. It is difficult to quantify a worst case acceptable sampling range across a bit's period.

Usually, RS232 communications are performed over very short distances; regarding RS485 as an empirical law between baud rate and cable length may help:

$$\text{length(m)} = \frac{1200}{\sqrt{1 + \left(\frac{\text{baudrate}}{250}\right)^2}}$$

Baud rate is expressed in Kbit/s.

## 4.5 Timing limit of UART communications

This section investigates the maximum allowed deviation of the baud rates of sender and receiver for correct communication.

The single bit cells of a UART frame are analyzed and the maximum admissible delay until bits are detected incorrectly is calculated. Sender and receiver use baud rates that deviate from the ideal baud rate. The worst case is a slow receiver trying to listen to a sender transmitting too fast. The other extreme with a slow transmitter and a fast receiver is less critical. For correct transmission, the instant of the detection point of the last bit cell of a frame must be prior the instant of the end of the transmitted UART frame.

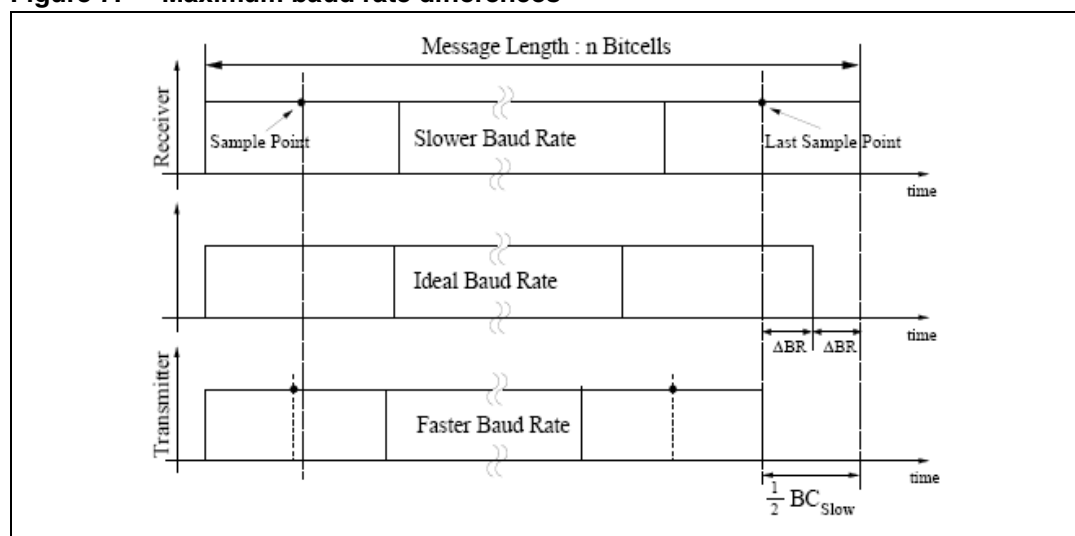
The condition for a successful communication is given in the following equation:

$$\frac{n - \frac{1}{2}}{\text{Baudrate}_s} - \frac{n}{\text{Baudrate}_f} \leq 0$$

$n$  is the UART frame length in bit cells including start, parity and stop bits. By rewriting the last equation it follows that the ratio between Baud Rate<sub>fast</sub> and Baud Rate<sub>slow</sub> must fulfil a condition in order to guarantee correct communication:

$$\frac{\text{Baudrate}_f}{\text{Baudrate}_s} \leq \frac{n}{n - \frac{1}{2}}$$

For example, considering an 8N1 (8 data bits, no parity, 1 stop bit) frame, it follows that the Baud Rate<sub>fast</sub> can be up to 5.26% higher than Baud Rate<sub>slow</sub> in order to communicate correctly.

**Figure 7. Maximum baud rate differences**

## 4.6 Examples

As a first example, we consider a frame 8N1 received by an ST10F27x connected with a PC, baud rate is 115200 bit/s, the connection is made through a simple straight cable 2 meters long, the connectors are standard DB9 and an ST232 is used as level translator.

In this case, reading [Table 2 on page 9](#), if we choose a reload value of 0011h and  $SxBRS = 0$ , the arithmetic error is -3.55%. In this condition, the PC UART is faster than the ST10 UART. The transfers from the PC to the ST10 correspond to the worst case, being a fast transmitter and a slow receiver. The maximum receive jitter is  $1/(16 * 115200)$  seconds or in terms of percentage of the frame about 0.6%. If the reference clock is obtained through a standard ceramic resonator, we can assume a further 0.5% error whereas the PLL jitter is negligible. Finally, supposing that the capacitive effect of the cable reduces the percentage of the bit time available for sampling of about the 5% (very worst case), another 0.5% should be added to our computation.

In [Section 4.5](#) it has been shown that an 8N1 frame allows a difference of 5.26% between the baud rate of the fast transmitter and the slow receiver in order to communicate correctly. Even if all the factors described above sum up in a very unlucky case, a remaining safety margin of 0.1% allows an error-free communication.

As another example, we consider an 8E1 frame (8 data bits, even parity, 1 stop bit) sent by ST10F27x on an RS485 network 100 meters long (twisted pair cable), baud rate is 2 Mbit/s and an ST485 is the transceiver.

In this case there is no arithmetic error but the effects of the cable and the clock issues like the jitter have more impact. It's difficult to quantify those effects; in any case, the clock tolerance of 4.76% (see [Section 4.5](#)) also allows error-free communications with significative capacitive cable effects and not very stable clock references.

### Conclusion

Keeping in mind all the possible sources of problems and checking all the points described, correct UART communications can be achieved in most cases.

## 5 Revision history

**Table 6. Document revision history**

Date	Revision	Changes
27-Apr-2006	1	Initial release

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